

DAVE ROWLAND

WHAT CAN C++ LEARN ABOUT
THREAD SAFETY FROM OTHER
LANGUAGES?























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> An artist's conception of a supernova explosion. Credit: NASA's Goddard Space Flight Center / ESA / Hubble / L. Calcada



Borrowing Trouble: The Difficulties Of A C++ Borrow-Checker



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Introduction

A common question raised when comparing C++ and Rust is whether the Rust borrow checker is really unique to Rust, or if it can be implemented in C++ too. C++ is a very flexible language, so it seems like it should be possible. In this article we'll explore if it is possible to do borrow checking at compile time in C++.

Some background on C++ efforts

Many folks are working on <u>improving C++</u>, including improving its memory safety. <u>Clang</u> has <u>experimental -Wlifetime warnings</u> to help catch a class of use-after-free bugs. The cases it catches are typically <u>dangling references to temporaries</u>, which makes them a valuable set of warnings to enable when it is available. But the cases it would solve do not seem to intersect with the set of cases <u>MiraclePtr</u> is attempting to protect against, which is an effort to frustrate

Merging state and references breaks ownership

If we accept that we can modify the language to make HasMut<T> and HasRef<T> nondestructible, and to enforce they are not used after a move, then we might consider to go a step further and do away with these troublesome types.

We might try to instead make the reference types MutRef<T> and Ref<T> not-publiclydestructible but also movable with a destructive move. Then we can eliminate the HasMut and HasRef types, and encode those states by the existence of the reference types.

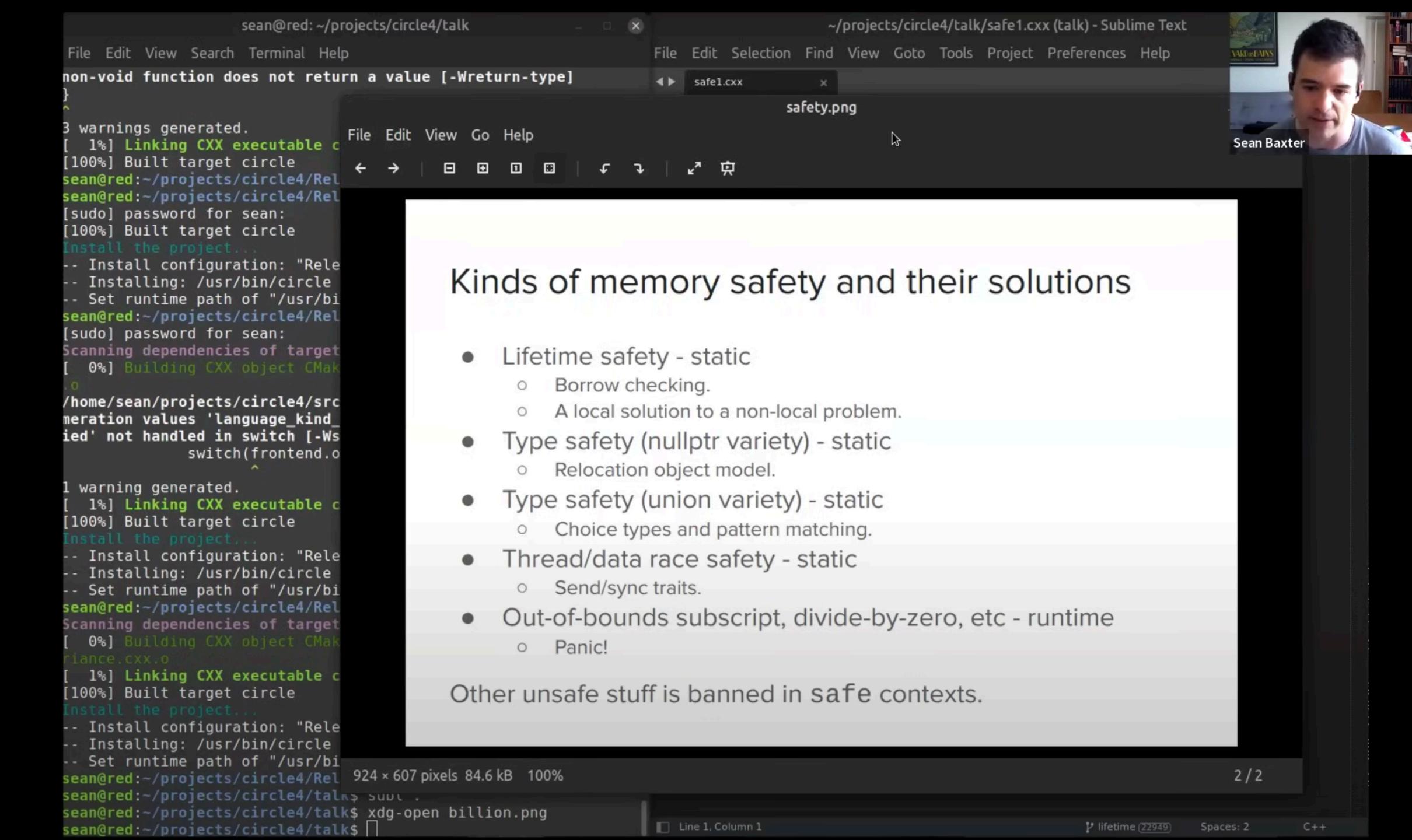
However, that allows a method to steal ownership from a reference. By constructing a Uniq<T> from a MutRef<T>, ownership is taken without being passed a Uniq<T> explicitly. Thus we actually need the states representing HasMut and HasRef to remain in the original scope of the Uniq<T> they are transitioned from in order to return ownership back to the same scope (though not the same variable).

Conclusion

We attempted to represent ownership and borrowing through the C++ type system, however the language does not lend itself to this. Thus memory safety in C++ would need to be achieved through runtime checks.



"However, the language does not lend itself to this. Thus memory safety in C++ would need to be achieved through runtime checks."

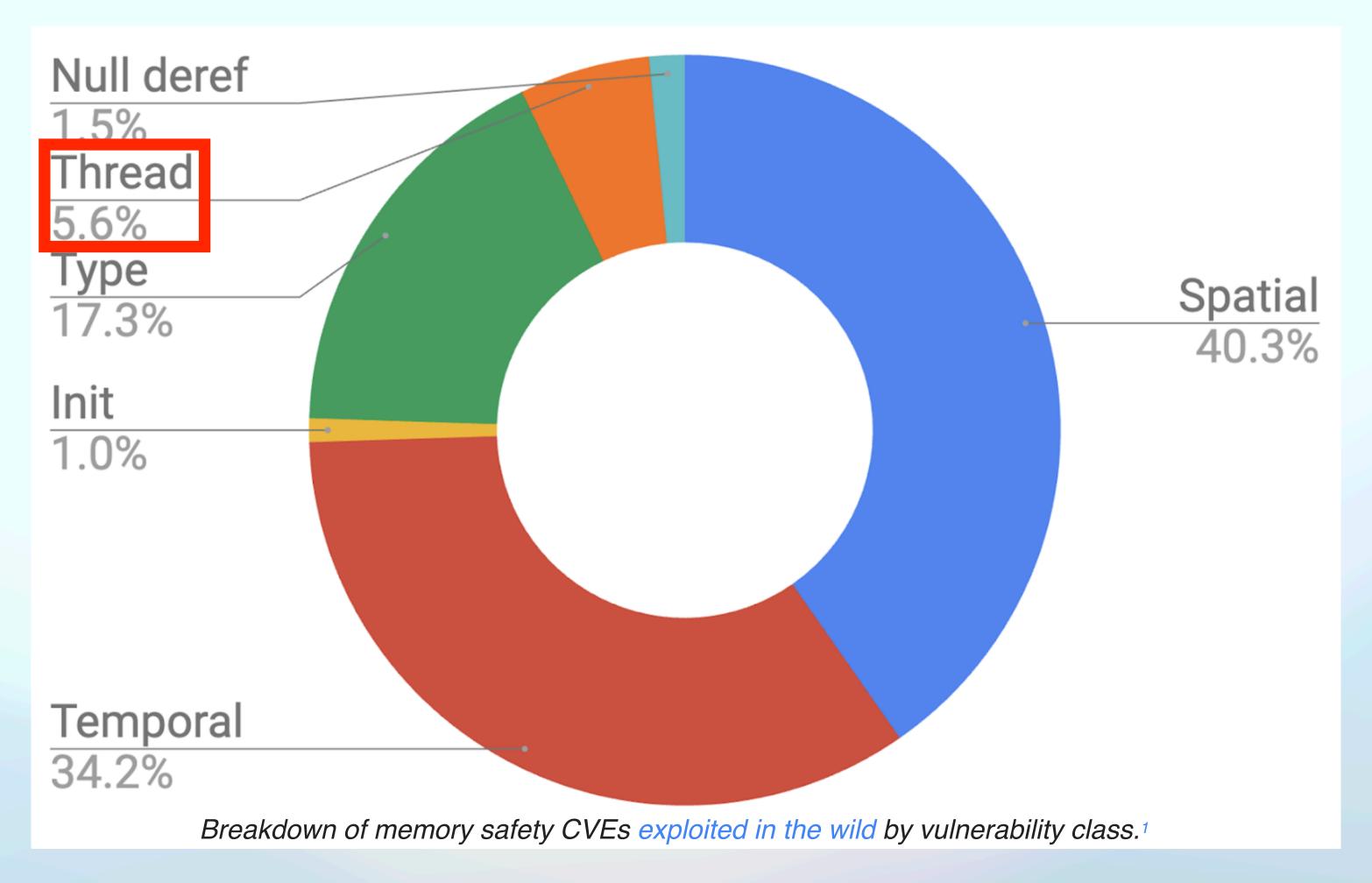


Thread safety:

What is Thread Safety?

What is Thread Safety?

- A program is thread safe if it is free from data races
 - And dead-locks
- A data race is when two threads access the same memory location when at least one of them is a write
- A thread safe programming language makes it impossible to express data races



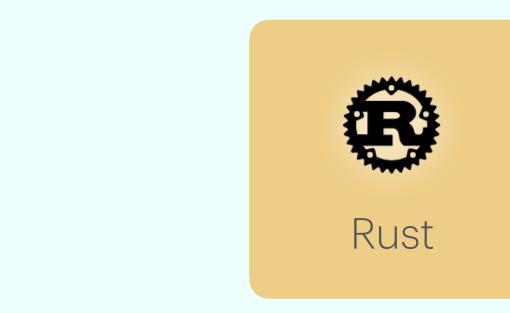
Source: Retrofitting spatial safety to hundreds of millions of lines of C++ https://security.googleblog.com/2024/11/retrofitting-spatial-safety-to-hundreds.html

- 5.6% seems small but will grow
 - As "low hanging fruit" memory safety improves
 - As machines gain more cores
 - As multi-threading becomes easier and more ubiquitous e.g. std::execution
 - We will see lifetime/temporal safety is inextricably to thread safety

- Bugs are difficult to spot and difficult to debug
- Problems typically arise long after a data race occurs
- Even if only **5.6%** of bugs are thread related, the time spent fixing them is likely much higher

- 5.6% is in Google's research
- In some industries that are inherently real-time (like audio) this is likely to be much higher

Landscape of Approaches





Haskell









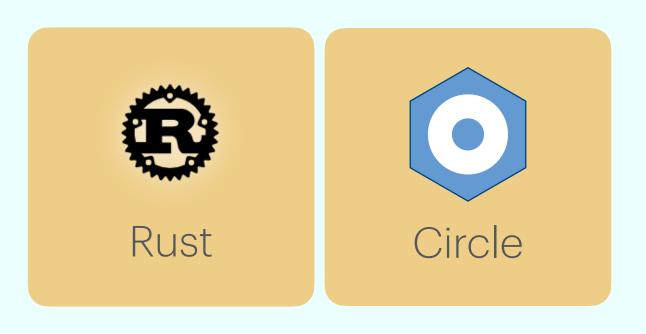


Immutability
(Pure value semantics)

Mutable value semantics (Local reasoning)

Message passing (No shared state)

Focus On





Sync & Send

Actors

Low-level

High-level









Sync & Send

- Protocols/traits that are checked
- A sync object can be safely shared between threads
- A send object can be safely transferred between threads

Sync & Send in Swift

The Sendable Protocol

- Notion of "isolation boundaries" between potential thread execution contexts
- · Objects can only pass isolation boundaries if they conform to the @Sendable protocol
 - Sendable can be inferred in some cases
- Syncable objects are a special case of Sendable objects
 - E.g. a LockingResource
 - No "syncable" keyword

```
open class Thread : NSObject {
    public convenience init(block: @escaping @Sendable () -> Void)
```

Sync & Send in Rust

- send is a "marker trait"
 - Similar to a C++ "type trait"
- Inferred if:
 - A copy can be made (value semantics)
 - A borrow can shared (T&)
 - NOT mutable borrow (mut T&)

Sync & Send in Circle

- send is a "marker interface"
 - Similar to a C++ "type trait"
- Inferred if:
 - A copy can be made (value semantics)
 - A borrow can shared (const T^)
 - NOT mutable borrow (T^)

```
sean@red: ~/projects/circle4/talk
File Edit View Search Terminal Help
an owned place is a local variable or subobject of a local variab
g is a non-local variable declared at rell.cxx:8:6
Pair g { 10, 20 };
sean@red:~/projects/circle4/talk$ circle match1.cxx
      match1.cxx:21:10
 return match(obj) {
match-expression is not exhaustive
 .i8, .u8, .i16, .u16, .u32, .i64, .s
sean@red:~/projects/circle4/talk$ circle thread1.cxx
rror: thread1.cxx:22:32
      threads^.push back(thread(&entry point, ^s, i));
error during overload resolution for std2::thread::thread
 instantiation: std2.h:1225:9
   thread/(where F:static, Args...:static)(F f, Args... args) sa
 during constraints checking of template parameter Args
 template arguments: [
    F = void(&)(std2::basic_string<char, std2::allocator<char>>^/
SCC-0, int) safe
   Args#0 = std2::basic_string<char, std2::allocator<char>>^/
   Args#1 = int
   constraint: std2.h:1224:26
      template<std2::send F, std2::send... Args>
   constraint std2::send not satisfied over std2::basic_string<c
har, std2::allocator<char>>⁴
sean@red:~/projects/circle4/talk$
```

```
File Edit Selection Find View Goto Tools Project Preferance
■ | | | | r match1.cxx x | match2.cxx x | match3.cxx x | std2.h x
     #feature on safety
     #include "std2.h"
                                            Sean Baxter
     using namespace std2;
     // Can we pass mutable borrows into thread entry
     void entry point(string^ s, int tid) safe {
       s^->append("More text");
       // println(*s);
 10
 11
     int main() safe {
       vector<thread> threads { };
 13
 14
 15
 16
         // s dies before the threads join, so possib
 17
         string s = "Hello threads";
 18
 19
         // Launch all threads.
          const int num threads = 15;
 20
21
          for(int i : num threads)
            threads^.push back(thread(&entry point, ^s.
22
 23
24
 25
       // Join all threads.
       for(thread^ t : ^threads)
26
27
         t^->join();
 28
```

~/projects/circle4/talk/thread1.cxx (talk) - Sublime



Sync and send in C++?

scl - Safe Concurrency Library



```
class safe_thread
public:
    template<typename F, send... Args>
    safe_thread (F&& f, Args&&... args)
        : thread (std::forward<F> (f), std::forward<Args> (args)...)
       // N.B. We can't constrain F to the concept due to recursion of is move constructable
        // So we have to statically assert it
        static_assert (send<F>);
    safe_thread (safe_thread&& other)
        : thread (std::move (other.thread))
private:
   std::jthread thread;
```



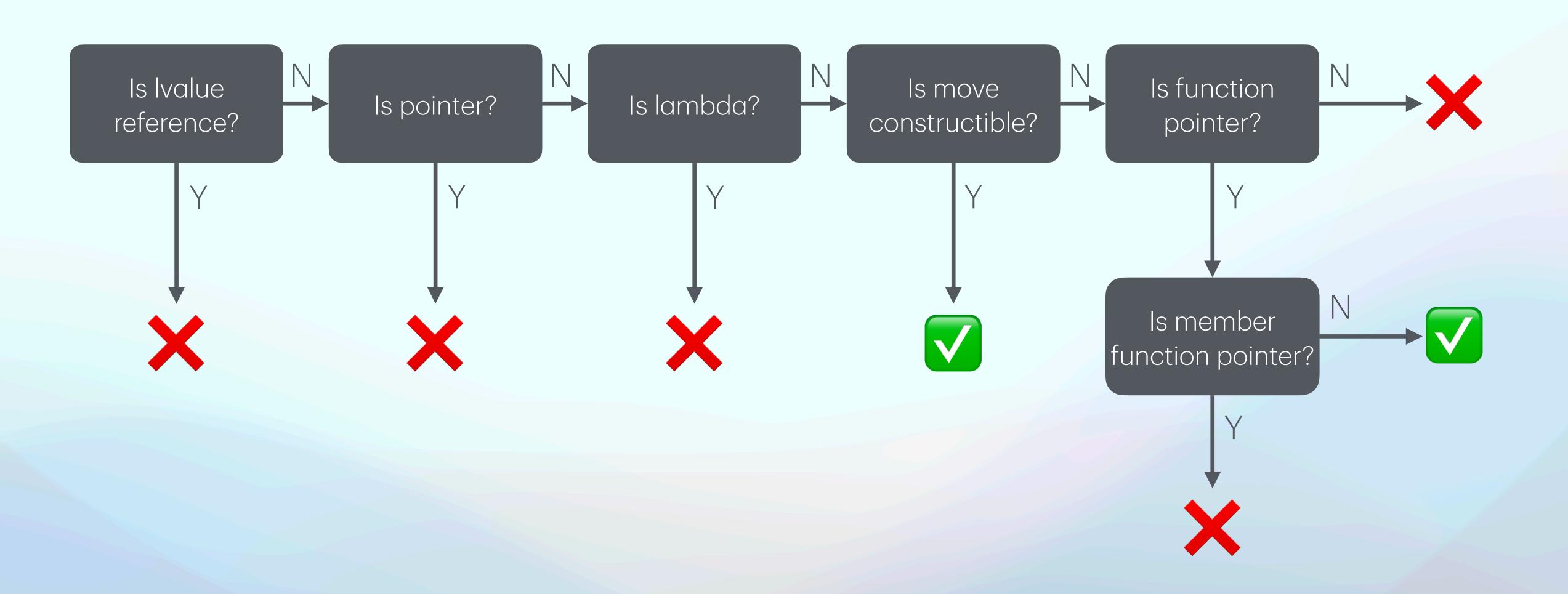
Send in C++: Moved between threads

```
template<typename F, send... Args>
safe thread (F&& f, Args&&... args)
    : thread (std::forward<F> (f), std::forward<Args> (args)...)
    static_assert (send<F>);
```

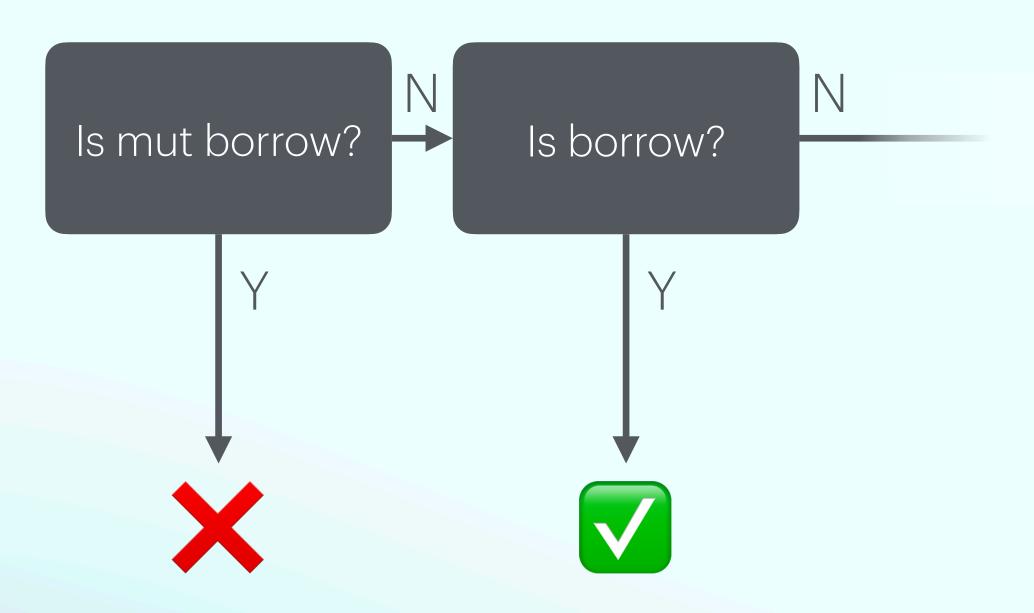
```
template<typename T>
struct is send : std::bool constant<</pre>
       (! (std::is_lvalue_reference_v<T>)
               std::is_pointer_v<std::remove_extent_t<T>>
               is lambda v<T>))
         &&
           (std::is_move_constructible_v<T>
            || (is_function_pointer_v<std::decay_t<T>>
                && ! std::is_member_function_pointer_v<T>))>
{};
template<typename T>
concept send = is_send<T>::value;
```

```
static_assert(is send v<const int>);
static assert(is send v<int>);
static_assert(is send v<int&&>);
static assert(is send v<int>);
static_assert(! is send v<int&>);
static_assert(! is send v<int*&>);
static_assert(! is send v<const int&>);
static_assert(! is send v<const int*&>);
static_assert(! is send v<std::string&>);
static_assert(! is send v<const std::string&>);
static assert(! is send v<std::string*&>);
static_assert(! is_send_v<const std::string*&>);
```

Is T send?



R Is T send?



Send in C++: Moved between threads

No

- Ivalue references
- Object pointers
- Lambdas
- May be referenced outside this thread boundary

Only

- rvalues
- Non-member function pointers
- Can be sure no data is shared

Sync in C++: Sharable between threads

```
template<typename T>
struct is_sync : std::false_type {};

template<typename T>
struct is_sync<std::atomic<T>> : std::true_type {};

template<typename T>
inline constexpr bool is_sync_v = is_sync<T>::value;

template<typename... Args>
concept sync = (is_sync<Args>::value && ...);
```

```
static_assert(! is_sync_v<int>);
static_assert(! is_sync_v<int&>);
static_assert(! is_sync_v<const int&>);
static_assert(! is_sync_v<std::string&>);
static_assert(! is_sync_v<const std::string&>);
static_assert(! is_sync_v<std::atomic<int>>);
```

What types are sync?

- std::
 - std::atomic
 - Trivial types only
- synchronized_value (P0290)
 - Wraps a type with a std::mutex
 - Automatically locks during access
 - Works with any type

synchronized_value

```
template<typename Type>
class synchronized value
public:
    synchronized value(const synchronized value&) = delete;
    synchronized value &operator=(const synchronized value&) = delete;
    template<typename... Args>
    synchronized value(Args&&... args)
        : val (std::forward<Args> (args)...)
    template<typename Fn, typename Up, typename... Types>
    friend std::invoke_result_t<Fn, Up&, Types&...> apply (Fn&&, synchronized_value<Up>&,
                                                            synchronized value<Types>&...);
private:
    std::mutex mutex;
    Type val;
                               template<typename T>
};
                               struct is sync<synchronized value<T>> : std::true type
                               {};
```

```
template<sync T>
struct is_send<std::shared_ptr<T>> : std::true_type
{};
```

Good

```
std::shared_ptr<std::atomic<int>>
```

std::shared_ptr<synchronized_value<std::string>>



Bad

```
std::shared_ptr<int>
```

std::shared_ptr<std::string>



```
void entry_point (std::shared_ptr<synchronized_value<std::string>> sync_s
                                                                           int tid)
   apply ([tid] (auto& s) {
        s.append ("%");
       std::println ("{} {}", s, tid);
       return s;
    *sync_s);
int main()
   auto s = std::make shared<synchronized value<std::string>> ("Hello threads");
    std::vector<safe_thread> threads { };
    const int num threads = 15;
    for (int i : std::views::iota (0, num threads))
       threads.push_back | safe_thread (entry_point, auto (s), auto (i)));
```



Problems: Nested Pointers

```
struct node
   node* next;
   node* prev;
```

```
void entry point (std::shared ptr<synchronized value<<del>std::string></del>> sync s, int tid)
    apply ([tid] (auto& s) {
        //...
        return s;
    },
    *sync s);
int main()
        //...
        auto s = std::make shared<synchronized value<<del>std::string></del>> ("Hello threads");
        //...
```

Problems: this Pointers

```
threads.push_back (safe_thread (entry_point, auto (s), auto (i)));
```

Problems: Global Pointers

```
void set global string (std::string*);
void entry point (std::shared ptr<synchronized value<std::string>> sync s, int tid)
    apply ([tid] (auto& s) {
        set global string (&s);
       //...
       return s;
    },
    *sync s);
int main()
       //...
        auto s = std::make shared<synchronized value<std::string>> ("Hello threads");
       //...
```

Problems: Leaked Pointers

```
auto widget = std::make_unique<Widget> (args);
auto widget ptr = widget.get();
threads.push_back (safe_thread (entry_point, std::move (widget)));
widget_ptr->do_stuff();
```

Problems: Summary

- Nested pointers
- this pointers
- Global pointers
- Leaked pointers

How far have we got in C++?

Safer, but not safeTM

How far have we got in C++?

- Used an unenforceable safe_thread class
- Used a non-standard synchronized_value class
 - Had to add our own type trait for it
- Did a lot of fighting with the compiler
 - Template instantiation
 - Similar to "fighting the borrow checker"?
- Added a lot of overhead to our code
 - Atomic reference counting
 - Mutex locking

How far have we got in C++?

- Not bullet proof
- Not beginner friendly
- Not default





```
void entry_point (std::shared_ptr<synchronized_value<std::string>> sync_s, int tid)
{
    apply ([tid] (auto& s) {
        s.append ("\( \beta \)");
        std::println ("{} {}", s, tid);
        return s;
    },
    *sync_s);
}
int main()
{
    auto s = std::make_shared<synchronized_value<std::string>> ("Hello threads");
    std::vector<safe_thread> threads { };
    const int num_threads = 15;
    for (int i : std::views::iota (0, num_threads))
        threads.push_back (safe_thread (entry_point, auto (s), auto (i)));
}
```





```
void entry_point (
       std::shared_ptr<synchronized_value<std::string>> data,
       int tid)
   apply ([tid] (auto& s) {
       s.append ("");
       std::println ("{} {}", s, tid);
        return s;
   *data);
int main()
   //...
   threads.push_back (safe_thread (entry point.
                                    auto (s), auto (i)));
```

```
void entry_point (
        shared_ptr<mutex<string>> data,
        int thread_id) safe
    auto lock_guard = data->lock();
    string^s = lock_guard^.borrow();
    s^->append ("");
int main() safe
    //...
    threads _ push_back(thread (&entry point.
                               copy shared_data, i));
```

Same example in Rust

```
use std::sync::{Arc, Mutex};
use std::thread;
fn entry_point(data: Arc<Mutex<String>>, thread_id: i32) {
   let mut guard = data.lock().unwrap();
    guard.push_str("\(\beta\)");
   println!("Thread {}: {}", thread_id, *guard);
pub fn main() {
   let shared_data = Arc::new(Mutex::new(String::from("Hello threads")));
   let mut threads = Vec::new();
    const NUM_THREADS: i32 = 15;
   for i in 0..NUM_THREADS {
        // Clone the Arc for this thread
        let data_clone = Arc::clone(&shared_data);
        // Spawn the thread and store its handle
        let handle = thread::spawn(move | | {
            entry_point(data_clone, i);
        });
        threads.push(handle);
    for handle in threads {
        handle.join().unwrap();
```

Same example in Rust (with borrows)

```
use std::sync::Mutex;
use std::thread;
fn entry_point(data: &Mutex<String>, thread_id: i32) {
    let mut guard = data.lock().unwrap();
    guard.push str("\(\beta\)");
    println!("Thread {}: {}", thread_id, *guard);
pub fn main() {
   let shared data = Mutex::new(String::from("Hello threads"));
    const NUM THREADS: i32 = 15;
    // Use scope to ensure threads don't outlive our data
    thread::scope(|scope| {
        let mut threads = Vec::new();
        for i in 0..NUM THREADS {
            let local_data = &shared_data;
            let handle = scope.spawn(move | | {
                entry_point(local_data, i);
            });
            threads.push(handle);
        for handle in threads {
            handle.join().unwrap();
    });
```

Key changes made in this version:

- Removed Arc and now using direct references (&Mutex<String>)
- 2. Added thread::scope to ensure threads don't outlive the borrowed data
- 3. Changed the thread spawning to use scoped threads via scope spawn
- 4. Simplified the function signature of entry_point to take a reference
- 5. No more need for explicit cloning since we're using references

Same example in Rust (with borrows)

```
use std::sync::Mutex;
use std::thread;
fn entry_point(data: &Mutex<String>, thread_id: i32) {
    let mut guard = data.lock().unwrap();
    guard.push str("\(\beta\)");
    println!("Thread {}: {}", thread_id, *guard);
pub fn main() {
   let shared data = Mutex::new(String::from("Hello threads"));
    const NUM_THREADS: i32 = 15;
    // Use scope to ensure threads don't outlive our data
    thread::scope(|scope| {
        let mut threads = Vec::new();
        for i in 0..NUM THREADS {
            let local_data = &shared_data;
            let handle = scope.spawn(move | | {
                entry_point(local_data, i);
            });
            threads.push(handle);
        for handle in threads {
            handle.join().unwrap();
    });
```

This version has several advantages:

- More efficient (no atomic reference counting)
- Cleaner code (no clone operations)
- Compile-time guarantees about data lifetime
- Still maintains thread safety through the Mutex

Without a way to properly express lifetimes (in terms of borrows/relocations/drops) we don't get the same level of safety and performance

Back to C++

```
void entry_point (std::shared_ptr<synchronized_value<std::string>> sync s, int tid)
    apply ([tid] (auto& s) {
        s.append ("%");
        std::println ("{} {}", s, tid);
        return s;
   *sync s);
int main()
    auto s = std::make shared<synchronized value<std::string>> ("Hello threads");
    std::vector<safe thread> threads { };
   const int num threads = 15;
    for (int i : std::views::iota (0, num_threads))
        threads.push back (safe thread (entry point, auto (s), auto (i)));
```

Problems: Summary

- Nested pointers
- this pointers
- Global pointers
- Leaked pointers

C++ Reflection to the Rescue?

Recursive Sync/Send Type Trait Checking

- Check members of types are all sendable
- Check members of lambdas are all sendable

```
struct node
{
    node* next;
    node* prev;
};

std::shared_ptr<syncronized_value<node>>();
```

```
consteval auto is_send_type (std::meta::info type) -> bool
   type = remove_cv (type);
   // Non-member function pointers
   if (is_pointer_type (type)
                                                                  co|{
       && is_function_type (remove_pointer (type))
       && ! is_member_function_pointer_type (type))
       return true;
   // lvalue refs and pointers
   if (is_lvalue_reference_type (type)
        || is_pointer_type (remove_extent (type)))
      return false;
   // POD built-in types
   if (is_arithmetic_type (type))
        return true;
   // Recursive class/struct/lambda members
   if (is_class_type (type))
        return std::ranges::all_of(nonstatic_data_members_of(type),
                                    [](std::meta::info d)
                                       return is_send_type (type_of(d));
                                   });
   // Construct from rvalue ref
   if (is_rvalue_reference_type (type)
       && is_constructible_type (type, { remove_reference (type) }))
       return true;
    return false;
```

```
template<typename T>
inline constaynr hool is send v = is send (^^T).
template<typename T>
consteval auto is_send() -> bool

if (is_send_type (^^T))
    return true;

return is_sync_v<T>;
}

template<typename T>
inline constexpr bool is_send_v = is_send<T>();
template<typename T>
concept send = is_send_v<T>;
```

```
[[maybe unused]] auto non capturing = [] (int) {};
struct node
                                         static_assert(is_send_v<decltype(non_capturing)>);
   node* prev;
                                         int i = 0:
   node* next;
                                         [[maybe_unused]] auto val_capturing = [i] (int) {};
};
                                         static_assert(is_send_v<decltype(val_capturing)>);
static_assert(! is_send_v<node>);
                                          [[maybe_unused]] auto ref_capturing = [&i] (int) {};
struct type
                                         static_assert(! is_send_v<decltype(ref_capturing)>);
    type()
        auto n = std::make_shared<node>();
        [[maybe_unused]] auto this_capturing = [this] { run(); };
        static_assert(! is_send_v<decltype(this_capturing)>);
        [[maybe_unused]] auto this_n_capturing = [this, n] { run(); };
        static_assert(! is_send_v<decltype(this_n_capturing)>);
        [[maybe_unused]] auto n_ref_capturing = [&n] {};
        static_assert(! is_send_v<decltype(n_ref_capturing)>);
        [[maybe_unused]] auto n_val_capturing = [n] {};
        static_assert(! is_send_v<decltype(n_val_capturing)>);
    void run() {}
```

|};

Problems: Summary

- Nested pointers
- this pointers
- Global pointers
- Leaked pointers

Global pointers

```
void set_global_string (std::string*);

void entry_point (std::shared_ptr<synchronized_value<std::string>> sync_s, int tid)
{
    apply ([tid] (auto& s) {
        set_global_string (&s);
        //...
        return s;
    },
    *sync_s);
}
```



```
fn entry_point(data: &Mutex<String>, thread_id: i32) {
   let mut guard = data.lock().unwrap();
   guard.push_str(".");
   println!("Thread {}: {}", thread_id, *guard);
}
```

Wrapping with Reflection

- P3294 Code Injection with Token Sequences
 Hopeful for C++29
- P0707 Metaclasses

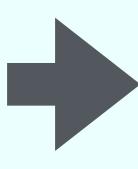
 Proposed !



Implicit synchronized_value

metaclass proposed syntax

```
class person
public:
    person() = default;
    std::string get_first_name() const
        return first_name;
    void set_first_name (std::string_view new_first)
        first_name = new_first;
    // Repeat for last_name
private:
    std::string first_name, last_name;
```



```
class person
public:
   person() = default;
    std::string get_first_name() const
      return apply ([] (auto& p) {
                         return p.get_first_name();
                     person_internal);
   void set_first_name (std::string_view new_first)
       apply ([&] (auto& p) {
                   p.set_first_name (new_first);
               person_internal);
   // Repeat for last_name
private:
   struct ___
             _person;
   mutable synchronized_value<__person> person_;
```



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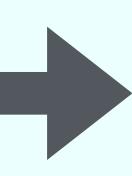


Now in EDG... godbolt.org/z/fex55qq50

```
consteval auto make_interface_functions(info proto) -> info {
   info ret = ^^{};
   for (info mem : members_of(proto)) {
       if (is_nonspecial_member_function(mem)) {
           ret = ^^{
               \tokens(ret)
               virtual [:\(return_type_of(mem)):]
                  \id(identifier_of(mem)) (\tokens(parameter_list_of(mem))) = 0;
           3;
          --- reporting compile time errors not yet implemented ---
       // else if (is_variable(mem)) {
            print
                  consteval void interface(std::meta::info proto) {
       // } // e
                       std::string_view name = identifier_of(proto);
   return ret;
                       queue_injection(^^{
                            class \id(name) {
                            public:
                                \tokens(make_interface_functions(proto))
                                virtual ~\id(name)() { }
                            3;
```

Implicit mutex locking

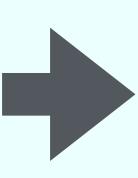
```
class person(mutex)
public:
    person() = default;
    std::string get_first_name() const
        return first_name;
    void set_first_name (std::string_view new_first)
        first_name = new_first;
    // Repeat for last_name
private:
    std::string first_name, last_name;
```



```
class person
public:
    person() = default;
    std::string get_first_name() const
       std::scoped_lock _ (mutex);
       return person__get_first_name();
    void set_first_name (std::string_view new_first)
        std::scoped_lock _ (mutex);
        person__set_tirst_name (new_first);
    // Repeat for last_name
private:
    class
            _person;
   std::mutex mutex;
   mutable __person person_;
};
template<>
struct is_sync<person> : std::true_type {};
```

Implicit shared_mutex locking

```
class person(shared_mutex)
public:
    person() = default;
    std::string get_first_name() const
        return first_name;
    void set_first_name (std::string_view new_first)
        first_name = new_first;
    // Repeat for last_name
private:
    std::string first_name, last_name;
```



```
class person
public:
    person() = default;
    std::string get_first_name() const
       std::shared_lock _ (mutex);
       return person__get_first_name();
    void set_first_name (std::string_view new_first)
       std::unique_lock _ (mutex);
        person_.set_tirst_name (new_first);
    // Repeat for last_name
private:
    class
            person;
   std::shared_mutex mutex;
    mutable __person person_;
};
template<>
struct is_sync<person> : std::true_type {};
```



```
void entry_point (std::shared_ptr<synchronized_value<std::string>> sync_s, int tid)
   apply ([tid] (auto& s) {
       s.append (".");
       std::println ("{} {}", s, tid);
       return s;
   *sync_s);
int main()
    auto p = std::make_shared<synchronized_value<std::string>> ("Hello threads");
    //...
```



```
void entry_point (std::shared_ptr<person> p, int tid)
   apply ([tid] (auto& s) {
       s.append ("");
       std::println ("{} {}", s, tid);
       return s;
   *sync_s);
int main()
   auto p = std::make_shared<person> ("Hello threads");
   //...
```



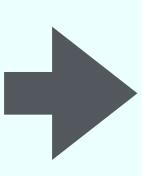
```
auto widget = std::make_unique<Widget> (args);
auto widget_ptr = widget.get();
threads.push_back (safe_thread (entry_point, std::move (widget)));
widget_ptr->do_stuff();
```

```
void entry_point (std::shared_ptr<person> p, int tid)
{
   auto person_ptr = p.get();
}
```

Wrapped std::shared_ptr

arc metaclass

```
class person(arc)
public:
    person() = default;
    std::string get_first_name() const
        return first_name;
    void set_first_name (std::string_view new_first)
        first_name = new_first;
    // Repeat for last_name
private:
    std::string first_name, last_name;
```



```
class person
public:
    person() = default;
    std::string get_first_name() const
       return person_->get_first_name();
    void set_first_name (std::string_view new_first)
       person_->set_first_name (new_first);
    // Repeat for last_name
private:
   class __person;
   std::shared_ptr<__person> person_;
```

```
class Person
    private var first_name: String = "";
    private var last_name: String = "";
    func get_first_name() -> String
        return first_name
    mutating func set_first_name (new_first: String)
        first_name = new_first;
    // Repeat for last_name
```

```
class person(arc)
public:
    std::string get_first_name() const
        return first_name;
    void set_first_name (std::string_view new_first)
        first_name = new_first;
    // Repeat for last_name
private:
    std::string first_name, last_name;
};
```



```
class person(mutex, arc)
{
public:
    //***
```



```
void entry_point (std::shared_ptr<person> p, int tid)
{
    p.set_first_name ("\black");
    std::println ("{} {}", p.get_first_name(), tid);
}
int main()
{
    auto p = std::make_shared<person> ("Hello threads");
    //...
}
```



```
void entry_point (person p, int tid)
{
    p.set_first_name ("**);
    std::println ("{} {}", p.get_first_name(), tid);
}
int main()
{
    auto p = person ("Hello threads");
    //...
}
```

Swift class: Breaking Cycles

- Cyclic references cause memory leaks
- References in Swift are strong by default
- To break a cycle **weak** references can be used
- These are niled when the last strong reference is destroyed
- Must be checked before dereferencing

```
var p = Person()
p.set_first_name (new_first: "Dave")
print (p.get_first_name())

weak var p2 = p
p2?.set_first_name (new_first: "John")
```

Wrapped std::weak_ptr

```
class person
public:
    class weak_ref
    public:
        weak_ref() = default;
        weak_ref (person p)
            : person_ (p.person_) {}
        std::optional<person> get() const
            if (auto valid = person_lock())
                return person (std::move (valid));
            return std::nullopt;
    private:
        std::weak_ptr<__person> person_;
    //... rest of class as before
private:
    person (std::shared_ptr<__person>&& other)
        : person_ (other) {}
};
```

```
person p1;
//... do stuff with p1
person::weak_ref p2; // create uninitialised
           // assign from strong-ref
p2 = p1;
if (auto valid_person = p3.get())
    std::println ("p3 {}", valid_person->get_first_name());
p2.get().transform ([] (auto valid_person) {
                       valid_person.set_first_name ("John");
                       return valid_person;
                   });
```

Swift structs

- Value semantics
 - Two instances of a **struct** have distinct objects
 - Are @Sendable implicitly if all members are @Sendable
 - Scalar (pod etc.) and self-contained objects (like String) are @Sendable
- Can be implemented with copy-on-write for efficiency

```
struct Person
   private var first_name: String =
   private var last_name: String = "";
   mutating func set_first_name (new_first: String)
       first_name = new_first;
    func get_first_name() -> String
        return first_name
   // Repeat for last_name
```

```
struct person
    std::string get_first_name() const
        return first_name;
    void set_first_name (std::string_view new_first)
        first_name = new_first;
    // Repeat for last_name
private:
    std::string first_name, last_name;
};
```



c.o.w. metaclass

```
struct person(cow)
    std::string get_first_name() const
        return first_name;
    void set_first_name (std::string_view)
        first_name = new_first;
   // Repeat for last_name
private:
    std::string first_name, last_name;
};
```

```
struct person
    std::string get first name() const {
        return person_->get_first_name();
    void set first name (std::string view new first) {
        copy if shared();
        person_->set_first_name (new_first);
    // Repeat for last name
private:
    struct
             person;
    static_assert (std::is_copy_constructible_v<__person>);
    std::shared_ptr<__person> person_
        = std::make_shared<__person>();
    void copy if shared() {
        if (person_.use_count() > 1)
            person = std::make shared< person> (*person );
```

Copy on Write structs

```
struct person
    std::string get first name() const {
        return person_->get_first_name();
   void set_first_name (std::string_view new_first) {
        copy if shared();
        person ->set first_name (new_first);
    // Repeat for last name
private:
    struct __person;
    static_assert (std::is_copy_constructible_v<__person>);
    std::shared ptr< person> person
        = std::make shared<__person>();
    void copy_if_shared() {
        if (person_.use_count() > 1)
            person_ = std::make_shared<__person> (*person_);
```

- Each person has its own shared_ptr instance
 - This is never shared
- As soon as a non-const function is called, a unique copy is made
 - The internal __person may be shared, but that's fine as there will only be readers



- Only works if there are no pointers or references to a person
 - send enforces this when passed to a thread
 - Delete operator new to avoid heap allocations
 - Delete operator& to avoid taking the address
- Doesn't stop references
- Doesn't stop references/pointers when used as a member in another object
 - Would require viral checking/static analysis

```
struct person
{
    //...
    // Wrapped __person functions
    //...
};
```

Mutable Value Semantics

- Hylo's thread safety comes from avoiding shared state
- Objects are mutable within a function local reasoning
- Similar to Swift with only struct types
- Implemented efficiently



- send trait introduces an "isolation boundary" between threads
 - Objects can only be copied or moved between them
- sync trait tells the compiler an object is data-race free
 - And is implicitly send
- These traits need to be checked recursively for all members
 - C++23 can not do this
 - C++26 reflection should enable this checking
- Lifetime safety is inherently intertwined with thread safety
 - Solved in other languages with borrow checking or mutable value semantics
- · We need to encapsulate pointers in value types to ensure they're not exposed to abuse
 - C++26 Reflection generation (and future metaclasses) can make this simple

Limitations

- Not the most efficient
 - E.g. mutex wraps the whole class, not individual members
- Can arc, cow or mutex metaclasses be inherited?
 - If any original functions were virtual, this would break protections e.g. cow, mutex
 - Derived classes could possibly inherit the metaclasses?
 - Could only work on non-virtual classes
 - Could be disabled by adding final to the generated class
- Very early days!
 - Need implementation experience

Concerns

- Bakes data-race safety and lifetime management in to the type
 - May not be suitable for every use case
 - Could pay performance cost for simple, single thread uses
 - Not the most efficient (borrow checking)
 - Great success in existing languages e.g. Swift
- Not "C++"?
 - Contradicts "Don't pay for what you don't use"

Sync & Send

Actors

Low-level

High-level









Actors

High-level



Actors

High-level



Swift Actors

```
actor Person
    private var first_name: String = "";
    func set_first_name (new_first: String)
        first_name = new_first;
    func get_first_name() -> String
        return first_name
```

```
var p = Person();
await p.set_first_name (new_first: "Dave")
print (await p.get_first_name())
```



metaclass proposed syntax

```
class person
public:
    person() = default;
    std::string get_first_name() const
        return first_name;
   void set_first_name (std::string new_first)
        first_name = new_first;
   // Repeat for last_name
private:
    std::string first_name, last_name;
```





```
auto get_scheduler()
{
    static exec::static_thread_pool pool(1);
    return pool.get_scheduler();
}
```

```
class person
public:
    std::string get_first_name() const
   void set_first_name (std::string new_first)
private:
   mutable __person person;
```





```
std::println ("\t\t\tmain tid: {}", std::this_thread::get_id());

person p;
std::println ("Name: {}", p.get_first_name());

std::thread t ([&]
{
    std::println ("\t\t\thread tid: {}", std::this_thread::get_id());

    p.set_first_name ("Dave");
    std::println ("Name: {}", p.get_first_name());
}
t.join();
```

```
main tid: 134711587358592
get tid: 134711584224832
Name:
thread tid: 126536174790208
set tid: 134711584224832
get tid: 134711584224832
Name: Dave
```





```
std::string first_name = co_await person.get_first_name();
```

```
std::string first_name = co_await person.get_first_name();
```





```
actor Person
{
    private var first_name: String = "";

    func set_first_name (n: String) {
        first_name = n;
    }

    func get_first_name() -> String {
        return first_name
    }
}
```

```
var p = Person();

await p.set_first_name (new_first: "Dave")
print (await p.get_first_name())
```

```
struct person(actor)
{
    std::string get_first_name() const {
        return first_name;
    }

    void set_first_name (std::string n) {
        first_name = n;
    }

private:
    std::string first_name;
};
```

```
person p;

co_await p.set_first_name ("Dave");
std::print (co_await p.get_first_name());
```

Actors: Problems

- Thread/Lifetime safety issues with function arguments
 - assert the arguments are send?
 - Reflect on the lambda type to ensure it's send?
 - Forward arguments like we did for safe_thread?



- In practice may need different pools
 - Serialises all actors on to a single thread
- Could use "Annotations for Reflection" P3394
 - Different pool tags
 - Different schedular types

```
struct [[=MainActor]] person(actor)
//...

template<typename PoolType>
auto get_main_scheduler()
{
    static exec::run_loop loop {};
    // Needs to be dispatched by main thread return loop.get_scheduler();
}
```

```
auto get_scheduler()
{
    static exec::static_thread_pool pool(1);
    return pool.get_scheduler();
}
```

```
struct [[=LowPriority]] person(actor)
//...

struct LowPriority_tag;

template<typename PoolType>
auto get_scheduler()
{
    static exec::static_thread_pool pool(1);
    // init low-priority
    return pool.get_scheduler();
}
```

Actors: Problems

- Huge overhead to queue every operation on a thread
- Re-entrant functions should execute synchronously

Run-time Data Race Detection

Existing Strategy: TSan

- Only available in clang and gcc (no Visual Studio support)
- Requires separate running
- Mutually exclusive with other sanitisers (ASan, UBSan etc.)
- Only as good as test coverage
 - Fuzzing can help

- Extremely heavyweight
 - 5-15x slower execution
 - 5-10x increase in memory usage
 - 2-3x increase in binary size
 - Moderate increase in compilation time

Lightweight Data Race Detection

	No Readers	Active	Active
	No Writers	Reader	Writer
Read Enter	No race	No race	DATA RACE
Write	No race	DATA	DATA
Enter		RACE	RACE

Lightweight Data Race Detection

```
void read_started (check_state& state)
   ++state num_readers; // must be first
   if (state is writing)
        std::terminate();
        // read during active write
void write_started (check_state& state)
   // must be first
   if (state.is_writing.exchange (true))
        std::terminate();
        // write during active write
   if (state num_readers > 0)
        std::terminate();
        // write during active read
```

```
struct check_state
    std::atomic<size_t> num_readers { 0 };
    std::atomic<bool> is_writing { false };
};
```

```
void read_ended (check_state& state)
   --state.num_readers;
void write_ended (check_state& state)
   state.is_writing = false;
```



Lightweight Data Race Detection

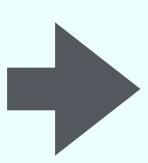
```
enum class check_type
    read,
   write
template<check_type type>
struct scoped_check
    scoped_check (check_state& check_state)
        : state (check_state)
        if constexpr (type == check_type::read)
            read_started (state);
        else
            write_started (state);
   ~scoped_check()
        if constexpr (type == check_type::read)
            read_ended (state);
        else
            write_ended (state);
    check_state& state;
};
```



Wrapped Data Race Detection

data_race_checker metaclass

```
struct person(data_race_checker)
    std::string get_first_name() const
        return first_name;
    void set_first_name (std::string_view new_first)
        first_name = new_first;
    // Repeat for last_name
private:
    std::string first_name, last_name;
```



```
struct person
   std::string get_first_name() const
       scoped_check<check_type::read> _ (check_state);
       return person_.get_tirst_name();
   void set_first_name (std::string_view new_first)
       scoped_check<check_type::write> _ (check_state);
       person_.set_tirst_name (new_tirst);
   // Repeat for last_name
private:
   struct __person;
     person person ;
   mutable check_state check_state;
```



16 Library introduction

16.4 Library-wide requirements

16.4.6 Conforming implementations

16.4.6.10 Data race avoidance

- This subclause specifies requirements that implementations shall meet to prevent data race function shall meet each requirement unless otherwise specified. Implementations may prother than those specified below.
- ² A C++ standard library function shall not directly or indirectly access objects ([intro.multithread]) accessible by threads other than the current thread unless the objects are accessed directly or indirectly via the function's arguments, including this.
- ³ A C++ standard library function shall not directly or indirectly modify objects ([intro.multithread]) accessible by threads other than the current thread unless the objects are accessed directly or indirectly via the function's non-const arguments, including this.
- 4 [Note 1: This means, for example, that implementations can't use an object with static storage duration for internal purposes without synchronization because doing so can cause a data race even in programs that do not explicitly share objects between threads. end note]
- ⁵ A C++ standard library function shall not access objects indirectly accessible via its arguments or via elements of its container arguments except by invoking functions required by its specification on those container elements.
- Operations on iterators obtained by calling a standard library container or string member function may access the underlying container, but shall not modify it.
 - [Note 2: In particular, container operations that invalidate iterators conflict with operations on iterators associated with that container. end note]
- Implementations may share their own internal objects between threads if the objects are not visible to users and are protected against data races.
- ⁸ Unless otherwise specified, C++ standard library functions shall perform all operations solely within the current thread if those operations have effects that are visible to users.
- 9 [Note 3: This allows implementations to parallelize operations if there are no visible side effects. end note]

23 Containers library

23.2 Requirements

[container.requirements]

[containers]

23.2.3 Container data races

[container.requirements.dataraces]

- For purposes of avoiding data races ([res.on.data.races]), implementations shall consider the following functions to be const: begin, end, rbegin, rend, front, back, data, find, lower_bound, upper_bound, equal_range, at and, except in associative or unordered associative containers, operator[].
- Notwithstanding [res.on.data.races], implementations are required to avoid data races when the contents of the contained object in different elements in the same container, excepting vector
bool>, are modified concurrently.
- [Note 1: For a vector<int> x with a size greater than one, x[1] = 5 and *x.begin() = 10 can be executed concurrently without a data race, but x[0] = 5 and *x.begin() = 10 executed concurrently can result in a data race. As an exception to the general rule, for a vector

 | y, y[0] = true can race with y[1] = true. end note

std Containers

- A C++ standard library function shall not directly or indirectly modify objects ([intro.multithread]) accessible by threads other than the current thread unless the objects are accessed directly or indirectly via the function's non-const arguments, including this.
- For purposes of avoiding data races ([res.on.data.races]), implementations shall consider the following functions to be const: begin, end, rbegin, rend, front, back, data, find, lower _bound, upper_bound, equal_range, at and, except in associative or unordered associative containers, operator[].



All **const** member functions can be called concurrently by different threads on the same container.





Herb Sutter on software development

My little New Year's Week project (and maybe one for you?)

[Updates: Clarified that an intrusive discriminator would be far beyond what most people mean by "C++ ABI break." Mentioned unique addresses and common initial sequences. Added "unknown" state for passing to opaque functions.]

Here is my little New Year's Week project: Trying to write a small library to enable compiler support for automatic raw union member access checking.

The problem, and what's needed

During 2024, I started thinking: What would it take to make C/C++ union accesses type-checked? Obviously, the ideal is to change naked union types to something safe.(*) But because it will take time and effort for the world to adopt any solution that requires making source code changes, I wondered how much of the safety we might be able to get, at what overhead cost, just by recompiling existing code in a way that instruments ordinary union objects?

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I'm an author and speaker, and a programming language nerd whose focus is on enabling our program code to be both clean and fast. I've been writing about programming since 1993, usually about C++ or about concurrency and parallelism. I'm the designer

Avoiding ABI Breaks: Extrinsic Storage

```
// That's it. Here's an example:
// {
    union Test { int a; double b; };

// Test t = {42};
    std::cout << t.a;
    t.b = 3.14159;
    std::cout << t.b;
    std::cout << t.b;
    // std::cout << t.b;
    // y
}</pre>

// Interior in the count in the c
```

Avoiding ABI Breaks: Extrinsic Storage

```
class data_race_registry {
    static inline auto tags
                               = extrinsic_storage<check_state>{};
public:
   static inline auto get_state(void* pobj) noexcept {
        return *tags.find_or_insert(pobj);
    static inline auto on_destroy(void* pobj) noexcept -> void {
        tags.erase(pobj);
```

```
constexpr const_reference operator[](size_type __pos) const noexcept {
  _LIBCPP_ASSERT_VALID_ELEMENT_ACCESS(__pos <= size(), "string index out of bounds");
  scoped_check<check_type::read> _ (data_race_registry::get_state (this));
 if (__builtin_constant_p(__pos) && !__fits_in_sso(__pos))
   return *(__get_long_pointer() + __pos);
 return *(data() + __pos);
```



Data Races as Contract Violations

```
void read_started (check_state& state)
   ++state num_readers;
    contract_assert (! state.is_writing); // read during active write
void write_started (check_state& state)
    contract_assert (! state.is_writing.exchange (true)) // write during active write
    contract_assert (state.num_readers == 0) // write during active read
```

```
constexpr const_reference operator[](size_type __pos) const noexcept {
  _LIBCPP_ASSERT_VALID_ELEMENT_ACCESS(__pos <= size(), "string index out of bounds");
 scoped_check<check_type::read> _ (data_race_registry::get_state (this));
 if (__builtin_constant_p(__pos) && !__fits_in_sso(__pos))
   return *(__get_long_pointer() + __pos);
 return *(data() + __pos);
```

```
constexpr const_reference operator[](size_type ___pos) const noexcept
    pre (can_read(data_race_registry::get_state (this)))
{
    //...
```

```
basic_string& replace(size_type __pos1, size_type __n1, const basic_string& __str)
    pre (can_write(data_race_registry::get_state (this)))
{
    //...
```

Data Race Detection

- Extremely limited
 - Works on function entry/exit, not memory
 - All bets are off if functions return references/pointers
 - Only works on types that don't expose their memory
 - Member function delegation not shown
- Could be used to check container contracts
- Use TSan!



3.3. Profile: Concurrency

- **Definition**: no data races. No deadlocks. No races for external resources (e.g., for opening a file).
- Question: should we also deal with priority inversion, delays caused by excess contention on a lock? Suggested initial answer: no.
- **Observation**: The concurrency profile is currently the least mature of the suggested profiles. It has received essentially no work specifically related to profiles, but concurrency problems have received intensive scrutiny in other contexts (including the Core Guidelines and MISRA++) so I can offer a few suggestions for initial work:
 - o *Threads*: prefer **jthread** to **thread** to get fewer scope-related problems.
 - Dangling pointers: consider a jthread a container and apply the usual rules for resource lifetime (RAII) and invalidation (§3.9).
 - Aliasing: statically detect if a pointer is passed to another thread. For an initial version, that will require restrictions on pointer manipulation in non-trivial control flows. In general, not all aliasing can be detected statically, and we need to reject too complex code. Defining "too complex" is essential, or we will suffer portability problems because of compiler incompatibilities. See "Flow analysis" (§4).
 - Invalidation: use unique_ptr and containers without invalidation (e.g., gsl::dyn_array)
 to pass information between threads.
 - Mutability: Prefer to pass (and keep) pointers to const.
 - Synchronization: use scoped_lock to lessen the chance of deadlock. Look into the possibility of statically detecting aliases in more than one thread to mutable data and enforce the use of synchronization on access through them. Use unique_ptr combined with protecting against aliasing across treads.



Look into the possibility of statically detecting aliases in more than one thread to mutable data



3.3. Profile: Concurrency

- **Definition**: no data races. No deadlocks. No races for external resources (e.g., for opening a file).
- Question: should we also deal with priority inversion, delays caused by excess contention on a lock? Suggested initial answer: no.
- **Observation**: The concurrency profile is currently the least mature of the suggested profiles. It has received essentially no work specifically related to profiles, but concurrency problems have received intensive scrutiny in other contexts (including the Core Guidelines and MISRA++) so I can offer a few suggestions for initial work:
 - o *Threads*: prefer **jthread** to **thread** to get fewer scope-related problems.
 - Dangling pointers: consider a jthread a container and apply the usual rules for resource lifetime (RAII) and invalidation (§3.9).
 - Aliasing: statically detect if a pointer is passed to another thread. For an initial version, that will require restrictions on pointer manipulation in non-trivial control flows. In general, not all aliasing can be detected statically, and we need to reject too complex code. Defining "too complex" is essential, or we will suffer portability problems because of compiler incompatibilities. See "Flow analysis" (§4).
 - Invalidation: use unique_ptr and containers without invalidation (e.g., gsl::dyn_array)
 to pass information between threads.
 - Mutability: Prefer to pass (and keep) pointers to const.
 - Synchronization: use scoped_lock to lessen the chance of deadlock. Look into the possibility of statically detecting aliases in more than one thread to mutable data and enforce the use of synchronization on access through them. Use unique_ptr combined with protecting against aliasing across treads.



Mutability: Prefer to pass (and keep) pointers to const.



3.3. Profile: Concurrency

- **Definition**: no data races. No deadlocks. No races for external resources (e.g., for opening a file).
- Question: should we also deal with priority inversion, delays caused by excess contention on a lock? Suggested initial answer: no.
- **Observation**: The concurrency profile is currently the least mature of the suggested profiles. It has received essentially no work specifically related to profiles, but concurrency problems have received intensive scrutiny in other contexts (including the Core Guidelines and MISRA++) so I can offer a few suggestions for initial work:
 - o *Threads*: prefer **jthread** to **thread** to get fewer scope-related problems.
 - Dangling pointers: consider a jthread a container and apply the usual rules for resource lifetime (RAII) and invalidation (§3.9).
 - Aliasing: statically detect if a pointer is passed to another thread. For an initial version, that will require restrictions on pointer manipulation in non-trivial control flows. In general, not all aliasing can be detected statically, and we need to reject too complex code. Defining "too complex" is essential, or we will suffer portability problems because of compiler incompatibilities. See "Flow analysis" (§4).
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Aliasing: statically detect if a pointer is passed to another thread. <snip>

Conclusion

- C++ needs a way to identify "isolation boundaries"
 - l.e. send
- This introduces strong aliasing and lifetime requirements
- This is not compatible with existing pointers/references
- Reflection can help us write in the styles of other languages which have better thread safety
 - Safely encapsulates pointers
- For "C++ performance" and "Don't pay for what you don't use" we need borrow checking:
 - Sean Baxter: "Safe C++" wg21.link/P3390

What Can C++ Learn About Thread Safety From Other Languages

David Rowland



Questions?

Slides/video:

drowaudio.github.io/presentations

